

On the largest prime factor of $x^2 - 1$

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Abstract

In this paper, we find all integers x such that $x^2 - 1$ has only prime factors smaller than 100. This gives some interesting numerical corollaries. For example, for any positive integer n we can find the largest positive integer x such that all prime factors of each of $x, x + 1, \dots, x + n$ are less than 100.

Keywords Pell equation, Compact representation, Lucas sequence.

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1 Introduction

For any integer n we let $P(n)$ be the largest prime factor of n with the convention $P(0) = P(\pm 1) = 1$. Our main result in this paper is the determination of all integers x satisfying the inequality

$$P(x^2 - 1) < 100. \quad (1)$$

We also give some interesting corollaries to our main result. For simplicity, we will consider only solutions with positive values of x .

Before stating the results, some history. In 1964, Lehmer [6] found all positive integer solutions x to the inequality $P(x(x + 1)) \leq 41$. Notice that this amounts to finding all odd positive integers $y = 2x + 1$ such that $P(y^2 - 1) \leq 41$. There are 869 such solutions. In [7], the first author found all positive integer solutions of the similar looking equation $P(x^2 + 1) < 100$. There are 156 of them.

In both [6] and [7], the method of attack on this question was the following. Assume that x is a positive integers such that $P(x^2 \pm 1) \leq K$, where $K = 41$ if the sign is $-$ and $K = 100$ if the sign is $+$. Then we can write

$$x^2 \pm 1 = dy^2, \quad (2)$$

where d is squarefree, and $P(dy) \leq K$. When the sign is $-$, then there are 13 possible primes $p \leq 41$ which can participate in the factorization of d . When the sign is $+$, then $p \leq 100$ but either $p = 2$ or p is a prime which is congruent to 1 modulo 4. There is a total of 12 such primes. In both cases, we can write equation (2) in the form

$$x^2 - dy^2 = \mp 1.$$

Thus, our possible values for x appear as the first coordinate of one of the solutions of at most $2^{13} - 1 = 8191$ Pell equations. For a given Pell equation, the sequence $(y_n/y_1)_{n \geq 1}$ forms a Lucas sequence with real roots. The *Primitive Divisor Theorem* for Lucas sequences with real roots (see, for example, [3], or the more general result from [1] which applies to all Lucas sequences) says that if $n > 6$, then y_n has a prime factor which is at least as large as $n - 1$. In particular, for the cases treated in [6] (and [7]), it suffices to check the first 42 (respectively 100) values of the component x of the Pell equations involved there and among these one finds all possible solutions of the equations considered there.

We follow the same approach in the present paper. The first step is, as previously, to determine the form of all the possible solutions. This is done as explained above via the Primitive Divisor Theorem for the second component of the solutions to Pell equations. Since there are 25 primes $p < 100$, this leads to first solving $2^{25} - 1$ Pell equations $x^2 - dy^2 = 1$, the largest one of them having d with 37 decimal digits. This is clearly impossible using standard algorithms like the continued fractions, as the computations would be too slow and the fundamental solutions of some of the involved Pell equations would have hundreds of millions of decimal digits. Instead, the Pell equations are solved by first computing the regulator of the ring of integers of the corresponding quadratic field, and then from the regulator obtaining a compact representation of the fundamental solution of the Pell equation. The only algorithm known fast enough to compute the huge amount of regulators needed is Buchmann's subexponential algorithm. The output of this algorithm gives exactly the regulator under the Generalized Riemann Hypothesis and unconditionally is only a multiple of the regulator. Next, we check which of the solutions to the Pell equations will lead to solutions to equation (1). Finally, a check is performed that proves our search misses no solutions, thus eliminating the apparent dependence of our results on the Generalized Riemann Hypothesis.

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2 Application of the Primitive Divisor Theorem

Here, we explain in more detail the applicability of the Primitive Divisor Theorem alluded to in Section 1 to our problem.

Let x be an integer such that $x^2 - 1$ is a product of only the primes up to 97. We can then write $x^2 - 1 = dy^2$, or, equivalently, $x^2 - dy^2 = 1$, where

$$d = 2^{a_1} \cdot 3^{a_2} \cdots 97^{a_{25}}, \quad a_i \in \{0, 1\} \text{ for } i = 1, \dots, 25.$$

The only restriction for d above is that not all the a_i s can be zero. This is a Pell equation, so $(x, y) = (x_n, y_n)$, where $x_n + y_n\sqrt{d} = (x_1 + y_1\sqrt{d})^n$ for some positive integer n and $x_1 + y_1\sqrt{d}$ is the fundamental solution. As we have $2^{25} - 1$ possibilities for d , to get a finite number of possible solutions for our equation, we need an upper bound for n .

Let d be fixed, $\eta = x_1 + y_1\sqrt{d}$ and $\zeta = x_1 - y_1\sqrt{d}$. It can be easily seen that

$$\begin{cases} x_n + y_n\sqrt{d} &= (x_1 + y_1\sqrt{d})^n = \eta^n, \\ x_n - y_n\sqrt{d} &= (x_1 - y_1\sqrt{d})^n = \zeta^n. \end{cases}$$

From here, we deduce that

$$x_n = \frac{\eta^n + \zeta^n}{2} \quad \text{and} \quad y_n = \frac{\eta^n - \zeta^n}{2\sqrt{d}}.$$

It is easy to see that y_1 divides y_n . We define $u_n = y_n/y_1$. As $u_n = (\eta^n - \zeta^n)/(\eta - \zeta)$, it follows that $(u_n)_{n \geq 0}$ is a Lucas sequence of the first kind with real roots η and ζ . By a result of Carmichael (see [3]) known as the Primitive Divisor Theorem, it follows that if $n > 12$, then u_n has a *primitive divisor* p . This primitive divisor has several particular properties, the most important one for us being that it satisfies the congruence $p \equiv \pm 1 \pmod{n}$. This implies that for $n > 98$, there exists a prime $p \geq 101$ dividing u_n . Thus, $n \leq 98$.

3 The algorithm

We now give our algorithm. Let

$$S = \{2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, \dots, 97\},$$

be the set of all primes $p \leq 100$. The set S has 25 elements. Let R_d be the regulator of the ring of integers of the quadratic field $\mathbb{Q}(\sqrt{d})$ and let $x_1 + y_1\sqrt{d}$ be the fundamental solution of the Pell equation $x^2 - dy^2 = 1$.

For all $D \in \mathcal{P}(S)$

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1. $d = \prod_{p \in D} p$
2. Compute mR_d
3. Compute a compact representation of $x_m + y_m\sqrt{d}$
4. For $i = 1$ to 25 compute $\text{ord}_{p_i} y_m$
5. $z = 2^{\text{ord}_{p_1} y_m} \cdots 97^{\text{ord}_{p_{25}} y_m}$

6. Compute all the convergents p_n/q_n of the continued fraction expansion of \sqrt{d} having $q_n < z$, and check whether $p_n^2 - dq_n^2 = 1$

7. If $R_d - \log 2 - \log \sqrt{d} \approx \log z$

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$m = 1$ and x_1 is a solution

8. For $i = 2$ to 98

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For $j = 1$ to 25 compute $\text{ord}_{p_j} y_i$

$z = 2^{\text{ord}_{p_1} y_i} \cdot \dots \cdot 97^{\text{ord}_{p_{25}} y_i}$

If $i \cdot R_d - \log 2 - \log \sqrt{d} \approx \log z$, x_i is a solution

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The algorithm searches through all $2^{25} - 1$ possible d .

In step 1, a value for d is chosen.

In step 2, Buchmann's subexponential algorithm is used to compute R_d . This algorithm returns a multiple mR_d of the regulator R_d , unconditionally. If the Generalized Riemann Hypothesis is true, then $m = 1$. We will remove the dependence of our algorithm on the Generalized Riemann Hypothesis in step 6.

A *compact representation* of an algebraic number $\beta \in \mathbb{Q}(\sqrt{d})$ is a representation of β of the form

$$\beta = \prod_{j=1}^k \left(\frac{\alpha_j}{d_j} \right)^{2^{k-j}}, \quad (3)$$

where $d_j \in \mathbb{Z}$, $\alpha_j = (a_j + b_j \sqrt{d})/2 \in \mathbb{Q}(\sqrt{d})$, $a_j, b_j \in \mathbb{Z}$, $j = 1, \dots, k$, and k , α and d_j have $O(\log \sqrt{d})$ digits.

In step 3, a compact representation of $x_m + y_m \sqrt{d}$ is constructed using the algorithms described in [8]. The reason for using compact representations is that the standard representation of the fundamental solution of the Pell equation has $O(\sqrt{d})$ digits. Using the standard representation would make these computations impossible. More details about compact representations can be found in [5].

Since we only have the compact representation of the $x_m + y_m \sqrt{d}$, in step 4, we use algorithms from [9] to perform modular arithmetic on the compact representation. The p -adic valuations of y_m when p is one of the first 25 primes are computed in the following way. We define $v = \prod_{i=1}^{25} p_i$. We first compute $y_m \pmod{v}$. If $y_m \equiv 0 \pmod{v}$, we conclude that y_m is divisible by all 25 primes. Next, for all primes that we now know divide y_m , we compute $y_m \pmod{p_i^c}$, where c is a sufficiently large constant. We first take $c = 15$ and if $y_m \equiv 0 \pmod{p_i^c}$ still, we then replace c by $2c$ and repeat the computation.

In step 5, z , which is defined to be the part of y_m divisible by the first 25 primes, is computed.

The purpose of step 6 is to remove the dependence of this algorithm on the Generalized Riemann Hypothesis. If the Generalized Riemann Hypothesis is false, it is possible that, without this check, we could miss some solutions in our search. Suppose therefore that $m > 1$ and that y_k is such a solution that we are missing, meaning that y_k is divisible only by some of the first 25 primes for some k that is not a multiple of m . But $y_1 \mid y_k$, so y_1 is divisible only by some of the first 25 primes. Also $y_1 \mid y_m$, meaning that y_1 divides the part of y_m that is divisible by the first 25 primes, which in our notation is z . In other words, $y_1 \mid z$, implying $y_1 < z$. As y_1 has to be the denominator of a convergent of the continued fraction expansion of \sqrt{d} , it follows that by checking that the relation $p_n^2 - dq_n^2 \neq 1$ holds for all n satisfying $q_n < z$, we arrive at a contradiction. This proves that either $m = 1$, or that y_k has a prime factor larger than 100 for all positive integers k . This implies that our algorithm finds all solutions to equation (1) unconditionally, and if $y_m = z$, then $m = 1$. In practice, z will be a relatively small number, so usually only about 10 convergents will need to be computed.

Since $x_1 \approx y_1\sqrt{d}$, it follows that

$$R_d = \log(x_1 + y_1\sqrt{d}) \approx \log(2y_1\sqrt{d}) = \log 2 + \log \sqrt{d} + \log y_1,$$

so we can determine whether $y_1 = z$ by the test in step 7. In this test, we took that $a \approx b$ if $|a - b| < 0.5$. If $y_1 \neq z$, then $y_1 \geq 101 \cdot z$, so $\log y_1 > 4.61 + \log z$. This shows that great numerical precision is not needed here. Just in case, we used 10 digits of precision.

Step 8 checks whether any of the x_n are solutions for $n = 2$ to 98. With the purpose of speeding up the algorithm, the fundamental solution was not powered when computing $\text{ord}_{p_j} y_i$. Instead, we computed the fundamental solution modulo the appropriate integer and then powering modulo that integer. Also, usually not all n have to be checked. This is a consequence of the fact that if a prime p divides y_k , then it divides y_{lk} for every positive integer l . For example, if we get that y_2 is divisible by a prime larger than 100, then y_4, y_6, \dots do not need to be checked.

The running time of the algorithm is dominated by the computation of the regulator in step 2. Step 2 makes one the algorithm run in subexponential time. It is the only part of the algorithm that is not polynomial. The computations were performed on a Intel Xeon E5430. The computation of the regulators took around 12 days of CPU time, while the rest of the computations took around 3 days of CPU time.

Suppose that one wants to find the solutions of $P(x^2 - 1) < K$ using this algorithm. By the Prime Number Theorem, there are approximately $K/\log K$ primes up to K . This means that the algorithm will loop $2^{K \log K}$

times. The product of all these primes will be of size $O(e^K)$. This means that step 2 will run in $O(e^{\sqrt{K \log K}})$. Thus, the time complexity of the algorithm is $O(\exp(K/\log K + \sqrt{K \log K}))$.

4 Results

Theorem 1. a) *The largest three solutions x of the equation $P(x^2 - 1) < 100$ are*

$$x = \begin{cases} 19182937474703818751, \\ 332110803172167361, \\ 99913980938200001. \end{cases}$$

b) *The largest solution x of $P(x^4 - 1) < 100$ is $x = 4217$.*

c) *The largest solution x of $P(x^6 - 1) < 100$ is $x = 68$.*

d) *The largest n such that $P(x^{2n} - 1) < 100$ has an integer solution $x > 1$ is $n = 10$, the solution being $x = 2$.*

d) *The largest n such that $P(x^{2n} - 1) < 100$ has an integer solution $x > 2$ is $n = 6$, the solution being $x = 6$.*

e) *The equation $P(x^2 - 1) < 100$ has 16167 solutions.*

f) *The greatest power n of the fundamental solution of the Pell equation $(x_1 + y_1\sqrt{d})^n$ which leads to a solution of our problem is $(2 + \sqrt{3})^{18}$; i.e., $n = 18$ for $d = 3$. The case $d = 3$ also gives the most solutions, namely 10 of them.*

Proof:

The proof is achieved via a computer search using the algorithm from Section 3. Part b) is proved by finding the largest square of all the x , c) by finding the largest a cube, etc. \square

The largest solution x has 20 decimal digits, the second largest has 18 digits, followed by 5 solutions with 17 digits and 10 solutions with 16 digits. All of the mentioned large solutions are odd. The largest even solution x has 15 digits.

Theorem 2. *Write the equation (1) as*

$$x^2 - 1 = 2^{a_1} \dots 97^{a_{25}}.$$

Then the following hold:

a) *The solution with the largest number of $a_i \neq 0$ is $x = 9747977591754401$. For this solution, 17 of the a_i s are non-zero.*

- b) The solution with the largest $\sum_{i=1}^{25} a_i$ is $x = 19182937474703818751$.
For this solution, $\sum_{i=1}^{25} a_i = 47$.
- c) The single largest a_i appearing among all solutions is $a_1 = 27$ and corresponds to the solution $x = 4198129205249$.

Proof:

Again, this is done via a computer search using the algorithm from Section 3. \square

Dabrowski [4] considered a similar problem, where the prime factors of $x^2 - 1$ consist of the first k primes p_1, \dots, p_k . He formulated the following conjecture.

Conjecture 3. *The Diophantine equation*

$$x^2 - 1 = p_1^{\alpha_1} \cdots p_k^{\alpha_k}$$

has exactly 28 solutions $(x; \alpha_1, \dots, \alpha_k)$ in positive integers, as follows:

- a) (3; 3),
- b) (5; 3, 1), (7; 4, 1), (17; 5, 2),
- c) (11; 3, 1, 1), (19; 3, 2, 1), (31; 6, 1, 1), (49; 5, 1, 2), (161; 6, 4, 1),
- d) (29; 3, 1, 1, 1), (41; 4, 1, 1, 1), (71; 4, 2, 1, 1), (251; 3, 2, 3, 1), (449; 7, 2, 2, 1),
(4801; 7, 1, 2, 4), (8749; 3, 7, 4, 1),
- e) (769; 9, 1, 1, 1, 1), (881; 5, 2, 1, 2, 1), (1079; 4, 3, 1, 2, 1), (6049; 6, 3, 2, 1, 2),
(19601; 5, 4, 2, 2, 2),
- f) (3431; 4, 1, 1, 3, 1, 1), (4159; 7, 3, 1, 1, 1, 1), (246401; 8, 6, 2, 1, 1, 2),
- g) (1429; 3, 1, 1, 1, 1, 1, 1), (24751; 5, 2, 3, 1, 1, 1, 1), (388961; 6, 4, 1, 4, 1, 1, 1),
- h) (1267111; 4, 3, 1, 1, 3, 1, 1, 2).

The main result of [4] is that Conjecture 3 is true for $k \leq 5$. From our data, we confirm Dabrowski's conjecture in a much wider range.

Theorem 4. *Conjecture 3 is true for $k \leq 25$.*

Proof:

This is done by simply factoring all $x^2 - 1$, where x is a solution of (1). \square

The next theorem follows also trivially from our results. Recall that a positive integer n is K -smooth if $P(n) \leq K$. In particular, the main result of our paper is the determination of all the 100-smooth positive integers of the form $x^2 - 1$.

Corollary 5. *Let t be the largest odd solution, $t = 19182937474703818751$ and s be the largest even solution, $s = 473599589105798$ of equation (1).*

- a) *The largest consecutive 100-smooth integers are x and $x + 1$ where $x = (t - 1)/2$ is the largest solution of $P(x(x + 1)) < 100$.*
- b) *The largest consecutive even 100-smooth integers are $t - 1$ and $t + 1$.*
- c) *The largest consecutive odd 100-smooth integers are $s - 1$ and $s + 1$.*
- d) *The largest triangular 100-smooth integer is $(t^2 - 1)/8$.*

As we mentioned in the Introduction, the problem of finding two consecutive K -smooth integers was examined by Lehmer in [6] in the sixties. At that time, he was able to solve the above problem for the values $K \leq 41$. The advance of both computing power and theoretical arguments (namely, the compact representations) allow us to solve the much harder problem of finding consecutive K -smooth integers for any $K \leq 100$. Note that, as was already mentioned at the end of Section 3, the difficulty of this problem grows exponentially with K .

Our results can also be applied to finding k consecutive K -smooth integers, for any integer k . We obtain the following results.

Corollary 6. *The largest integer x satisfying*

$$P(x(x + 1) \dots (x + n)) < 100,$$

for a given n , are given in the following table:

n	x
1	9591468737351909375
2	407498958
3	97524
4	7565
5	7564
6	4896
7	4895
8	284

Proof:

To find k consecutive 100 smooth integers, we first create a list of all pairs of consecutive 100-smooth integers. Every odd solution of our starting problem will give us one such pair. This is because if x is an odd solution, then $(x - 1)/2$ and $(x + 1)/2$ are consecutive 100-smooth integers. Once this list is created, we search for overlaps in these pairs and obtain our results. \square

In the recent paper [10], Shorey and Tijdeman proved several extensions of some irreducibility theorems due to Schur. The main results of [10] rely heavily on Lehmer's results from [6]. Thus, replacing, for example Lemma 2.1 in [10] by our results, it is likely that the main results from [10] can be extended in a wider range.

Remark. The tables produced by our computations can be found on the webpage <http://web.math.hr/~fnajman>.

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